2023 Conference on Next Generation Arithmetic (CoNGA)

PLAUs: Posit Logarithmic Approximate Units to implement low-cost operations with real numbers

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Outline

Background

- Posit Logarithmic Approximate Units (PLAUs)
- Implementation analysis
- Applications
- Conclusions

Background

Scientific applications compute real numbers $(\log_2(x), \frac{1}{3})$ How to represent real numbers in computers?

- Floating-point (IEEE 754[™])
 - Used in most modern computers

1 bit	<i>n</i> bits	<i>m</i> bits
Sign	Exponent	Fraction

Emerging alternatives: bfloat16, TensorFloat, posits...

Posit[™] Arithmetic

- 1 bit New variable-length sign field: *Regime*
- Trade-off accuracy – dynamic range
- Only two special cases Zero and $\pm \infty$ (NaR)
- Single rounding mode
- Easy comparison



m bits

Fraction

Accuracy bits

k+1 bits

2 bits

Posit[™] Arithmetic

- New variable-length sign field: Regime
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Accuracy bits

Posit[™] Arithmetic

- New variable-length Sign Regime Exponent Fraction
 Fraction
- Trade-off accuracy – dynamic range
- Only two special cases Zero and $\pm \infty$ (NaR)
- Single rounding mode
- Easy comparison



Accuracy bits

Drawbacks?

Not implemented in computers

High Area/Latency overhead

Float

1 bit	n bits	<i>m</i> bits
Sign	Exponent	Fraction



Posit

1 bit	k+1 bi	ts	2 bits		<i>m</i> bits		
Sign	Regim	e Ex	ponent		Fraction		
1 bit		<i>k</i> +1 b	its		2 bits	<i>m</i> bits	
Sign		Regin	ne		Exponent	Fraction	
1 bit	k+1 bits	2 bits			<i>m</i> bits		_
Sign	Regime	Exponent		Fraction			-
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Posit Logarithmic Approximate Units (PLAUs)

Use logarithm properties

 $\log_b(xy) = \log_b(x) + \log_b(y)$

$$\log_b(\frac{x}{y}) = \log_b(x) - \log_b(y)$$

$$\log_b(\sqrt[y]{x}) = rac{\log_b(x)}{y}$$

and approximation

$$\log(1+x)pprox x$$
 , for small x



[1] John N. Mitchell, "Computer multiplication and division using binary logarithms." *IRE Transactions on Electronic Computers* 4 (1962): 512-517.

[2] Raul Murillo, et al. "PLAM: A posit logarithm-approximate multiplier." *IEEE Transactions on Emerging Topics in Computing* 10.4 (2021): 2079-2085.

Posit Logarithmic Approximation

Standard Posit value:

$$X = (-1)^s \times 2^{4r} \times 2^e \times (1+f)$$



$$\log_2 X = 4 \times r + e + \log_2(1+f)$$
$$\approx 4 \times r + e + f$$

 $f\in [0,1)$

PLAUs – Multiplication

k = 4 reg + exp

$$P_{exact} = (-1)^{s_A + s_B} \times 2^{k_A + k_B} \times (1 + f_A) \times (1 + f_B)$$

$$\begin{bmatrix} \log_2 X = \ \approx 4 \times r + e + f \end{bmatrix} \begin{bmatrix} \log_b(xy) = \log_b(x) + \log_b(y) \\ \log_b(xy) = \log_b(x) + \log_b(y) \end{bmatrix}$$
$$P_{approx} = \begin{cases} (-1)^{s_A + s_B} \times 2^{k_A + k_B} \times (1 + f_A + f_B) & \text{if } f_A + f_B < 1, \\ (-1)^{s_A + s_B} \times 2^{k_A + k_B + 1} \times (f_A + f_B) & \text{if } f_A + f_B \ge 1. \end{cases}$$

PLAUs – Division

$$k = 4 reg + exp$$

$$Q_{exact} = (-1)^{s_A - s_B} \times 2^{k_A - k_B} \times (1 + f_A)/(1 + f_B)$$

$$\log_2 X = \approx 4 \times r + e + f \qquad \log_b(\frac{x}{y}) = \log_b(x) - \log_b(y)$$

$$p_{approx} = \begin{cases} (-1)^{s_A - s_B} \times 2^{k_A - k_B} \times (1 + f_A - f_B) & \text{if } f_A - f_B \ge 0, \\ (-1)^{s_A - s_B} \times 2^{k_A - k_B - 1} \times (2 + f_A - f_B) & \text{if } f_A - f_B < 0. \end{cases}$$

PLAUs – Square Root

k = 4 reg + exp

$$R_{exact} = \sqrt{2^k \times (1+f)} = 2^{k/2} \times \sqrt{1+f}$$

$$\log_2 X = \approx 4 \times r + e + f \qquad \qquad \log_b(\sqrt[y]{x}) = \frac{\log_b(x)}{y}$$

$$R_{approx} = 2^{k/2} \times (1 + f/2)$$

Error for PLAUs



Error for PLAUs (II)



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Efficient HW Implementation for PLAUs



Efficient HW Implementation for PLAUs (II)

$$R_{approx} = 2^{k/2} \times (1 + f/2)$$

- Right shift
- Need to keep MSB
- Keep LSB for rounding



ASIC Synthesis – Multiplier



- +70% area savings
- 75% 80% less energy
- Up to 1150MHz

Exact Multiplier from:

[3] Raul Murillo, et al. "Comparing different decodings for posit arithmetic." *CoNGA 2022*.

CoNGA'23, Singapore, March 1-2, 2023

ASIC Synthesis – Divider

Synopsys DC 45-nm TSMC



- 81% 89% area savings
- 89% 91% power savings
- 99% 96% less energy
- Up to 1150MHz

ASIC Synthesis – Square Root

Synopsys DC 45-nm TSMC



- Area reduction by 82% – 91%
- ADP and energy savings of ~98% - 94%
- Up to 1150MHz

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Applications – Computer Vision

Blur filter (div)

Exact

Approximate





PSNR: 27.6019 SSIM: 0.98217





PSNR: 29.0023 SSIM: 0.98281





PSNR: 28.2892 SSIM: 0.98318





PSNR: 28.7468 SSIM: 0.97396

Applications – Computer Vision

Sobel filter for edge detection (mul, div, sqrt)

Exact

Approximate

PSNR: 51.0439 PSNR: 44.5427 PSNR: 45.6212 SSIM: 0.99909 SSIM: 0.99629 SSIM: 0.99629



PSNR: 38.2694 SSIM: 0.99503

Applications – Machine Learning

K-Nearest Neighbors (mul, sqrt)

Dataset	Instances	Attributes	Classes	K	Exact	Approx.
Iris	150	3	3	5	95.55%	95.55%
Wine	178	13	3	7	67.92%	67.92%
Glass	214	9	7	5	59.38%	59.38%
Breast Cancer	569	30	2	13	94.74%	94.74%

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Conclusions

- Posit arithmetic is raising interest for its properties...
- But still has a high implementation cost
- Posit Logarithmic Approximate Units (PLAUs)
 - Multiplication, division, and sqrt
 - Much faster circuits
 - Outstanding savings in area, power and energy
 - Relative error of -11.11%, 12.5% and 6.06%, respectively
 - Useful in error-tolerant applications (CV, ML, ...)

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Thank you!

PLAUs: Posit Logarithmic Approximate Units to implement low-cost operations with real numbers



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https://github.com/ RaulMurillo/CoNGA 23

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Backup slides

Example Posit16



$$p = ((1-3s)+f) imes 2^{(1-2s) imes (4r+e+s)}, \quad r = egin{cases} -k & ext{if } R_0 = 0 \ k-1 & ext{if } R_0 = 1 \end{cases}$$

$$p = ((1 - 3 \cdot 1) + rac{150}{2^8}) imes 2^{(1 - 2 \cdot 1) imes (4 \cdot 3 + 2 + 1)}, \quad r = 4 - 1 = 3$$

p = -0.000043154

Relative Error

$$E_{mul} = \begin{cases} \frac{1+f_A+f_B}{(1+f_A)(1+f_B)} - 1 & \text{if } f_A + f_B < 1, \\ \frac{2(f_A+f_B)}{(1+f_A)(1+f_B)} - 1 & \text{if } f_A + f_B \ge 1. \end{cases}$$
$$E_{div} = \begin{cases} \frac{(1+f_A-f_B)(1+f_B)}{(1+f_A)} - 1 & \text{if } f_A - f_B \ge 0, \\ \frac{(2+f_A-f_B)(1+f_B)}{2(1+f_A)} - 1 & \text{if } f_A - f_B < 0. \end{cases}$$
$$E_{sqrt} = \frac{1+f/2}{\sqrt{1+f}} - 1.$$







Fig. 5: Comparison of division implementations for Posit(32, 2).

Design.

- fixed

- ADDEDA

Design

- Esitt

MR Aves

Mill Paver

A07

P

BOS Energy

- Acpres

Comparison of Posit Unit and FPU

FPGA synthesis targeting Genesys II (Vivado 2020.2) [4]

Arithmetic unit (LUT, FF)	Posit(32,2)	+Single IEEE 754	+Double IEEE 754
PAU Area	(11879, 2985)	(11796, 2979)	(11810, 2979)
FPU Area	_	(3726, 1008)	(6352, 1905)
Total Area	(44693, 23636)	(50318, 25727)	(55900, 27652)

[4] D. Mallasén, R. Murillo, et al. "PERCIVAL: Open-Source Posit RISC-V Core with Quire Capability." arXiv preprint arXiv:2111.15286 (2021). Source code available at <u>github.com/artecs-group/PERCIVAL</u>